

Lecture 18: Pseudorandom Functions

Pseudo-random Functions (PRF)

- Let $\mathcal{G}_{m,n,k} = \{g_1, g_2, \dots, g_{2^k}\}$ be a set of functions such that each $g_i: \{0, 1\}^m \rightarrow \{0, 1\}^n$
- This set of functions $\mathcal{G}_{m,n,k}$ is called a pseudo-random function if the following holds.

Suppose we pick $g \xleftarrow{\$} \mathcal{G}_{m,n,k}$. Let $x_1, \dots, x_t \in \{0, 1\}^m$ be distinct inputs. Given $(x_1, g(x_1)), \dots, (x_{t-1}, g(x_{t-1}))$ for any computationally bounded party the value $g(x_t)$ appears to be uniformly random over $\{0, 1\}^n$

Secret-key Encryption using Pseudo-Random Functions

Before we construct a PRF, let us consider the following secret-key encryption scheme.

- 1 Gen(): Return $sk = id \xleftarrow{s} \{1, \dots, 2^k\}$
- 2 Enc_{id}(m): Pick a random $r \xleftarrow{s} \{0, 1\}^m$. Return $(m \oplus g_{id}(r), r)$, where $m \in \{0, 1\}^n$.
- 3 Dec_{id}(\tilde{c}, \tilde{r}): Return $\tilde{c} \oplus g_{id}(\tilde{r})$.

Features. Suppose the messages m_1, \dots, m_u are encrypted as the cipher-texts $(c_1, r_1), \dots, (c_u, r_u)$.

- As long as the r_1, \dots, r_u are all distinct, each one-time pad $g_{id}(r_1), \dots, g_{id}(r_u)$ appear uniform and independent of others to computationally bounded adversaries. So, this encryption scheme is secure against computationally bounded adversaries!
- The probability that any two of the randomness in r_1, \dots, r_u are not distinct is very small (We shall prove this later as “Birthday Paradox”)
- This scheme is a “state-less” encryption scheme. Alice and Bob do not need to remember any private state (except the secret-key sk)!

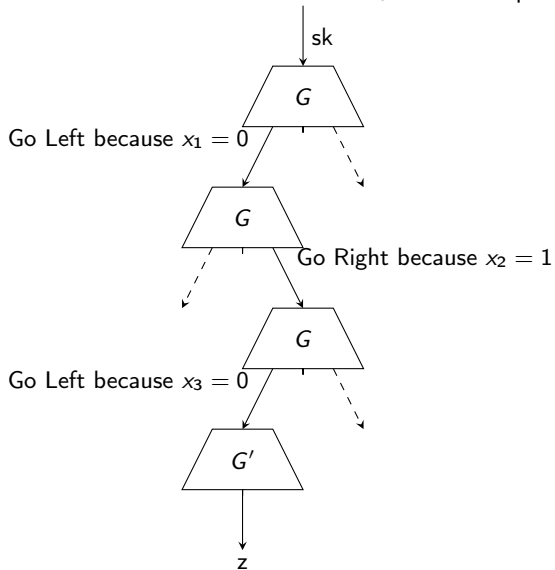
Construction of PRF I

- We shall consider the construction of Goldreich-Goldwasser-Micali (GGM) construction.
- Let $G: \{0, 1\}^k \rightarrow \{0, 1\}^{2k}$ be a PRG. We define $G(x) = (G_0(x), G_1(x))$, where $G_0, G_1: \{0, 1\}^k \rightarrow \{0, 1\}^k$
- Let $G': \{0, 1\}^k \rightarrow \{0, 1\}^n$ be a PRG
- We define $g_{\text{id}}(x_1 x_2 \dots x_m)$ as follows

$$G' (G_{x_m}(\dots G_{x_2}(G_{x_1}(\text{id}))\dots))$$

Construction of PRF II

Consider the execution for $x = x_1x_2x_3 = 010$. Output z is computed as follows.



We give the pseudocode of algorithms to construct PRG and PRF using a OWP $f: \{0, 1\}^{k/2} \rightarrow \{0, 1\}^{k/2}$

- Suppose $f: \{0, 1\}^{k/2} \rightarrow \{0, 1\}^{k/2}$ is a OWP
- We provide the pseudocode of a PRG $G: \{0, 1\}^k \rightarrow \{0, 1\}^t$, for any integer t , using the one-bit extension PRG construction of Goldreich-Levin hardcore predicate construction. Given input $s \in \{0, 1\}^k$, it outputs $G(s)$.

$G(k, t, s)$:

- 1 Interpret $s = (r, x)$, where $r, x \in \{0, 1\}^{k/2}$
- 2 Initialize bits = [] (i.e., an empty list)
- 3 Initialize $z = x$
- 4 For $i = 1$ to t :
 - 1 bits.append($\langle r, z \rangle$), here $\langle \cdot, \cdot \rangle$ is the inner-product
 - 2 $z = f(z)$
- 5 Return bits

- We provide the pseudocode of the PRF $g_{id}: \{0, 1\}^m \rightarrow \{0, 1\}^n$, where $id \in \{0, 1\}^k$, using the GGM construction. Given input $x \in \{0, 1\}^m$, it outputs $g_{id}(x)$.

$g(m, n, k, id, x)$:

- 1 Interpret $x = x_1x_2 \dots x_m$, where $x_1, \dots, x_m \in \{0, 1\}$
- 2 Initialize $inp = id$
- 3 For $i = 1$ to m :
 - 1 Let $y = G(k, 2k, inp)$
 - 2 If $x_i = 0$, then inp is the first k bits of y . Otherwise (if $x_i = 1$), inp is the last k bits of y .
- 4 Return $G(k, n, inp)$